1 Streams

In Python, we can use iterators to represent infinite sequences (for example, the generator for all natural numbers). However, Scheme does not support iterators. Let's see what happens when we try to use a Scheme list to represent an infinite sequence of natural numbers:

```
scm> (define (naturals n)
            (cons n (naturals (+ n 1))))
naturals
scm> (naturals 0)
Error: maximum recursion depth exceeded
```

Because **cons** is a regular procedure and both its operands must be evaluted before the pair is constructed, we cannot create an infinite sequence of integers using a Scheme list. Instead, our Scheme interpreter supports *streams*, which are *lazy* Scheme lists. The first element is represented explicitly, but the rest of the stream's elements are computed only when needed. Computing a value only when it's needed is also known as *lazy evaluation*.

```
scm> (define (naturals n)
                     (cons-stream n (naturals (+ n 1))))
naturals
scm> (define nat (naturals 0))
nat
scm> (car nat)
0
scm> (car nat)
#[promise (not forced)]
scm> (car (cdr-stream nat))
1
scm> (car (cdr-stream (cdr-stream nat)))
2
```

We use the special form cons-stream to create a stream: (cons-stream <operand1> <operand2>)

cons-stream is a special form because the second operand is not evaluated when evaluating the expression. To evaluate this expression, Scheme does the following:

- 1. Evaluate the first operand.
- 2. Construct a promise containing the second operand.
- 3. Return a pair containing the value of the first operand and the promise.

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To actually get the rest of the stream, we must call cdr-stream on it to force the promise to be evaluated. Note that this argument is only evaluated once and is then stored in the promise; subsequent calls to cdr-stream returns the value without recomputing it. This allows us to efficiently work with infinite streams like the naturals example above. We can see this in action by using a non-pure function to compute the rest of the stream:

```
scm> (define (compute-rest n)
...> (print 'evaluating!)
...> (cons-stream n nil))
compute-rest
scm> (define s (cons-stream 0 (compute-rest 1)))
s
scm> (car (cdr-stream s))
evaluating!
1
scm> (car (cdr-stream s))
1
```

Here, the expression compute-rest 1 is only evaluated the first time cons-stream is called, so the symbol evaluating! is only printed the first time.

When displaying a stream, the first element of the stream and the promise are displayed separated by a dot (this indicates that they are part of the same pair, with the promise as the cdr). If the value in the promise has not been evaluated by calling cdr-stream, we consider it to be not forced. Otherwise, we consider it forced.

```
scm> (define s (cons-stream 1 nil))
s
scm> s
(1 . #[promise (not forced)])
scm> (cdr-stream s) ; nil
()
scm> s
(1 . #[promise (forced)])
```

Streams are very similar to Scheme lists in that they are also recursive structures. Just like the cdr of a Scheme list is either another Scheme list or nil, the cdr-stream of a stream is either a stream or nil. The difference is that whereas both arguments to cons are evaluated upon calling cons, the second argument to cons-stream isn't evaluated until the first time that cdr-stream is called.

Here's a summary of what we just went over:

- nil is the empty stream
- cons-stream constructs a stream containing the value of the first operand and a promise to evaluate the second operand
- car returns the first element of the stream
- cdr-stream computes and returns the rest of stream

Questions

```
1.1 What would Scheme display?
    scm> (define (has-even? s)
           (cond ((null? s) #f)
                 ((even? (car s)) #t)
                 (else (has-even? (cdr-stream s)))))
   has-even?
    scm> (define (f x) (* 3 x))
    f
    scm> (define nums (cons-stream 1 (cons-stream (f 3) (cons-stream (f 5) nil))))
   nums
    scm> nums
    scm> (cdr nums)
    scm> (cdr-stream nums)
    scm> nums
    scm> (define (f x) (* 2 x))
    f
   scm> (cdr-stream nums)
    scm> (cdr-stream (cdr-stream nums))
    scm> (has-even? nums)
```

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1.2 Using streams can be tricky! Compare the following two implementations of filter-stream, the first is a correct implementation whereas the second is wrong in some way.

What's wrong with the second implementation?

1.3 Write a function map-stream, which takes a function f and a stream s. It returns a new stream which has all the elements from s, but with f applied to each one. (define (map-stream f s)

```
scm> (define evens (map-stream (lambda (x) (* x 2)) nat))
evens
scm> (car (cdr-stream evens))
2
```

1.4 Write a function slice which takes in a stream s, a start, and an end. It should return a Scheme list that contains the elements of s between index start and end, not including end. If the stream ends before end, you can return nil. (define (slice s start end)

```
scm> (define nat (naturals 0)) ; See naturals procedure defined earlier
nat
scm> (slice nat 4 12)
(4 5 6 7 8 9 10 11)
```

1.5 We can even represent the sequence of all prime numbers as an infinite stream! Define a function sieve, which takes in a stream of increasing numbers and returns a stream containing only those numbers which are not multiples of an earlier number in the stream. We can define primes by sifting all natural numbers starting at 2. Look online for the Sieve of Eratosthenes if you need some inspiration.

(define (sieve s)

(define primes
 (sieve (naturals 2)))
scm> (slice primes 0 10)
(2 3 5 7 11 13 17 19 23 29)

1.6 Since streams only evaluate the next element when they are needed, we can combine infinite streams together for interesting results! Use it to define a few of our favorite sequences. We've defined the function combine-with for you below, as well as an example of how to use it to define the stream of even numbers.

```
(define (combine-with f xs ys)
 (if (or (null? xs) (null? ys))
   nil
   (cons-stream
      (f (car xs) (car ys))
      (combine-with f (cdr-stream xs) (cdr-stream ys)))))
scm> (define evens (combine-with + (naturals 0) (naturals 0)))
evens
scm> (slice evens 0 10)
(0 2 4 6 8 10 12 14 16 18)
```

For these questions, you may use the naturals stream in addition to combine-with.

i. (define factorials

scm> (slice factorials 0 10)
(1 1 2 6 24 120 720 5040 40320 362880)

ii. (define fibs

scm> (slice fibs 0 10)
(0 1 1 2 3 5 8 13 21 34)

iii. (Extra for practice) Write exp, which returns a stream where the *n*th term represents the degree-*n* polynomial expansion for e^x , which is $\sum_{i=0}^{n} x^i/i!$.

You may use factorials in addition to combine-with and naturals in your solution. (define (exp x)

scm> (slice (exp 2) 0 5)
(1 3 5 6.333333333 7)

Extra Questions

1.1 In Discussion 9, we saw how to write a macro that creates a lambda function given an expression. While creating a parameter-less function might not seem that useful at first, it can be helpful in many cases when we don't want to immediately evaluate an expression.

Using the make-lambda macro you defined in Discussion 9, define make-stream, a macro which returns a pair of elements, where the second element is not evaluated until cdr-stream is called on it. Also define the procedure cdr-stream, which takes in a stream returned by make-stream and returns the result of evaluating the second element in the stream pair.

Unlike the streams we've seen in lecture and earlier in discussion, if you repeatedly call cdr-stream on a stream returned by make-stream, you may evaluate an expression multiple times.

```
(define-macro (make-stream first second)
```

```
(define (cdr-stream stream)
```

```
scm> (define a (make-stream (print 1) (make-stream (print 2) nil))))
1
a
scm> (define b (cdr-stream a))
2
b
scm> (cdr-stream b)
()
```